

B.Sc. III B.327 Linear Programming

Y2

- Q.1. In sensitivity analysis we deal with change in the optimal solution due to discrete variations in the parameters:
- c_i
 - a_{ij}
 - b_i
 - all of these
- Q.2 To maintain optimality of current optimal solution following change Δc_k is the coefficient c_k of nonbasic variable x_k , we have:
- $\Delta c_k = z_k - c_k$
 - $\Delta c_k \leq z_k - c_k$
 - $\Delta c_k > z_k - c_k$
 - none of these
- Q.3. Addition of a new constraint to the existing constraints of a L.P.P. will cause:
- a change in the objective function coefficients
 - no change in the objective function coefficients
 - a change in the existing optimal solution if it satisfies the new constraint
 - none of these
- Q.4. If the optimal solution to the L.P.P.
- $\text{Max } Z = -x_1 + 2x_2 - x_3$, subject to $3x_1 + 2x_2 - x_3 \leq 10$, $-x_1 + 4x_2 + x_3 \geq 6$ and $x_2 + x_3 \leq 4$, $x_1, x_2, x_3 \geq 0$, find the range of b_1 , constraint with optimal solution (given) $X_B = (6, 4, 10)$, $B^{-1} = \begin{bmatrix} 1 & 0 & -1 \\ 0 & 0 & 1 \\ 0 & -1 & 4 \end{bmatrix}$ and $x_1=0, x_2=4, x_3=0$ Max $Z=8$.
- $4 \leq b_1 < \infty$
 - $-6 \leq b_1 < \infty$
 - $-\infty < b_1 \leq 16$
 - none of these
- Q.5. If the optimal solution to L.P.P.
- $\text{Max } Z = 10x_1 + 3x_2 + 6x_3 + 5x_4$, subject to $x_1 + 2x_2 + x_4 \leq 6$, $3x_1 + 2x_3 \leq 5$, $x_2 + 4x_3 + 5x_4 \leq 3$ and $x_1, x_2, x_3, x_4 \geq 0$; if $x_1 = 5/3$, $x_2 = \frac{56}{27}$, $x_3 = 0$, $x_4 = 5/27$ then the limits for a_{23} so that new solution remains optimal feasible are (give that $a_{23} \notin B$, $\Delta_3 = -\frac{82}{27}$, $B_2 = \left[\frac{-5}{27}, \frac{1}{3}, \frac{1}{27} \right]$ and $C_B = 3$, $C_{B_2} = 10$, $C_{B_3} = 5$ in the last iteration)
- $\frac{39}{40} \leq a_{23} < \infty$
 - $-\frac{41}{10} \leq a_{23} < \infty$
 - $0 \leq a_{23} \leq \frac{81}{25}$
 - none of these

- Q.6. The advantages of dual simplex algorithm is that
 (a) It starts with a basic feasible solution (b) Artificial variable involved
 (c) No Artificial variable involved is it (d) None of these
- Q.7. If any variable of primal L.P.P. is unrestricted in sign, the corresponding constraint in dual will be
 (a) An equality (b) An inequality
 (c) Unrestricted in sign. (d) None of these
- Q.8. If the primal has an unbounded solution, the dual has either no solution or:
 (a) an unbounded solution (b) bounded solution
 (c) feasible solution (d) basic feasible solution
- Q.9. If the standard primal problem is of minimization, all the constraints involve the sign:
 (a) \geq (b) \leq (c) $=$ (d) None of these
- Q.10. The relation between the objective function of primal and dual problem of L.P.P. is
 (a) $\text{Max } Z_p = - \text{Min } Z_D$ (b) $\text{Max } Z_p = \text{Min } Z_D$
 (c) $\text{Min } Z_p = - \text{Max } Z_D$ (d) None of these
- Q.11. An integer programming problem can be solved by:
 (a) Gomory's cutting plane method (b) Hungarian method
 (c) Lowest cost entry method (d) None of these
- Q.12. Branch and Bound method to solve I.P.P. was developed by:
 (a) R.E. Gomory (b) Charnes and Cooper
 (c) A.H. Land and A.G. Doig (d) None of these
- Q.13. Gomory's cutting plane method may be used for solution of
 (a) all integer programming problem (b) mixed I.P.P.
 (c) both (a) & (b) (d) None of (a) & (b)
- Q.14. Branch and Bound technique may be used for solution of
 (a) all integer programming problem (b) mixed integer programming prob.
 (c) both (a) & (b) (d) None of (a) & (b)
- Q.15. For the solution of I.P.P. by Gomory method. The Gomory constraint equation is:
 (a) $\sum_{j \in R} f_{ij} x_j + x_{bi} = -f_{Bj}$ (b) $\sum_{j \in R} f'_{ij} x_j + x'_{bi} = -f'_{Bj}$
 (c) $\sum_{j \in R} d_{ij} x_j + x_{bi} = f'_{Bi}$ (d) None of these.